Advanced Systems Lab
Spring 2021

Lecture: SIMD extensions, AVX, compiler vectorization

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TA: Joao Rivera, several more

Flynn’s Taxonomy

<table>
<thead>
<tr>
<th></th>
<th>Single instruction</th>
<th>Multiple instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>Single data</td>
<td><strong>SISD</strong></td>
<td><strong>MISD</strong></td>
</tr>
<tr>
<td></td>
<td>Uniprocessor</td>
<td></td>
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<tr>
<td>Multiple data</td>
<td><strong>SIMD</strong></td>
<td><strong>MIMD</strong></td>
</tr>
<tr>
<td></td>
<td>Vector computer</td>
<td>Multiprocessors</td>
</tr>
<tr>
<td></td>
<td>Short vector extensions</td>
<td>VLIW</td>
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</tbody>
</table>
SIMD Extensions and AVX

AVX intrinsics

Compiler vectorization

The first version of this lecture (for SSE) was created together with Franz Franchetti (ECE, Carnegie Mellon) in 2008

Joao Rivera helped with the update to AVX in 2019

SIMD Vector Extensions

What is it?
- Extension of the ISA
- Data types and instructions for the parallel computation on short (length 2, 4, 8, ...) vectors of integers or floats
- Names: SSE, SSE2, AVX, AVX2 ...

Why do they exist?
- Useful: Many applications have the necessary fine-grain parallelism
  Then: speedup by a factor close to vector length
- Doable: Relatively easy to design by replicating functional units
Example AVX Family: Floating Point

Not drawn to scale

AVX: introduces three-operand instructions ($c = a + b$ vs. $a = a + b$)

AVX2: Introduces fused multiply-add (FMA: $c = c + a\times b$)

Sandy Bridge and later has (at least) AVX
Haswell/Skylake/ ...

Have AVX2

16 AVX registers

256 bit = 4 doubles = 8 singles

<table>
<thead>
<tr>
<th>%ymm0</th>
<th>%ymm8</th>
</tr>
</thead>
<tbody>
<tr>
<td>%ymm1</td>
<td>%ymm9</td>
</tr>
<tr>
<td>%ymm2</td>
<td>%ymm10</td>
</tr>
<tr>
<td>%ymm3</td>
<td>%ymm11</td>
</tr>
<tr>
<td>%ymm4</td>
<td>%ymm12</td>
</tr>
<tr>
<td>%ymm5</td>
<td>%ymm13</td>
</tr>
<tr>
<td>%ymm6</td>
<td>%ymm14</td>
</tr>
<tr>
<td>%ymm7</td>
<td>%ymm15</td>
</tr>
</tbody>
</table>

32 zmm (AVX-512)  16 ymm (AVX)  16 xmm (SSE)  scalar
AVX Registers

Used for different data types and instructions

256 bit

Integer vectors:
- 32-way byte
- 16-way 2 bytes
- 8-way 4 bytes
- 4-way 8 bytes

Floating point vectors:
- 8-way single
- 4-way double

Floating point scalars:
- single
- double

AVX Instructions: Examples

(three-operand!)

Double precision 4-way vector add: vaddpd %ymm1 %ymm0 %ymm1

(two-operand!)

Double precision scalar add (in SSE2): addsd %xmm0 %xmm1
Instruction Names (Assembly)

<table>
<thead>
<tr>
<th>AVX</th>
<th>SSE</th>
<th>Scalar (in SSE)</th>
</tr>
</thead>
<tbody>
<tr>
<td>packed (vector)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>vaddps</td>
<td>addps</td>
<td>addss</td>
</tr>
<tr>
<td>single precision</td>
<td></td>
<td></td>
</tr>
<tr>
<td>vaddpd</td>
<td>addpd</td>
<td>addsd</td>
</tr>
<tr>
<td>double precision</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Compiler will use this for floating point

x86-64 FP Code Example

Inner product of two vectors
- **Double precision arithmetic**
- **Compiled: not vectorized,** uses (single-slot) SSE instructions

```c
double ipf (double x[],
            double y[],
            int n) {
    int i;
    double result = 0.0;
    for (i = 0; i < n; i++)
        result += x[i]*y[i];
    return result;
}
```

```assembly
ifp:
    xorpd %xmm1, %xmm1
    xorl %ecx, %ecx
    jmp .L8
.L10:
    movslq %ecx,%rax
    incl %ecx
    movsd (%rsi,%rax,4), %xmm0
    mulsd (%rdi,%rax,4), %xmm0
    addsd %xmm0, %xmm1
.L8:
    cmpl %edx, %ecx
    jl .L10
    movapd %xmm1, %xmm0
    # result = 0.0
    # i = 0
    # goto middle
    # loop:
    # icpy = i
    # i++
    # t = y[icpy]
    # t *= x[icpy]
    # result += t
    # middle:
    # i:n
    # if < goto loop
    # return result
    ret
```
AVX: How to Take Advantage?

Necessary: fine grain parallelism

Options (ordered by effort):
- Use vectorized libraries (easy, not always available)
- Compiler vectorization (this lecture)
- Use intrinsics (this lecture)
- Write assembly

We will focus on floating point and double precision (4-way)

SIMD Extensions and AVX

Overview: AVX family

AVX intrinsics

Compiler vectorization

References:
Intel Intrinsics Guide
(easy access to all instructions, nicely done!)

Intel icc compiler manual
Visual Studio manual
Example AVX Family: Floating Point

Not drawn to scale

AVX: introduces three-operand instructions ($c = a + b$ vs. $a = a + b$)

AVX2: Introduces fused multiply-add (FMA)

Sandy Bridge and later has (at least) AVX

---

Intrinsics

Enable explicit use of vector instructions in C/C++

Assembly coded C functions

Expanded inline upon compilation: no overhead

Like writing assembly inside C

Floating point:
- Intrinsics for basic operations (add, mult, …)
- Intrinsics for math functions: log, sin, …

Our introduction is based on icc
- Almost all intrinsics work with gcc and Visual Studio (VS)
- Some language extensions are icc (or even VS) specific

Number of intrinsics

<table>
<thead>
<tr>
<th>ISA</th>
<th>Count</th>
</tr>
</thead>
<tbody>
<tr>
<td>MMX</td>
<td>124</td>
</tr>
<tr>
<td>SSE</td>
<td>154</td>
</tr>
<tr>
<td>SSE2</td>
<td>236</td>
</tr>
<tr>
<td>SSE3</td>
<td>11</td>
</tr>
<tr>
<td>SSSE3</td>
<td>32</td>
</tr>
<tr>
<td>SSE41</td>
<td>61</td>
</tr>
<tr>
<td>SSE42</td>
<td>19</td>
</tr>
<tr>
<td>AVX</td>
<td>188</td>
</tr>
<tr>
<td>AVX2</td>
<td>191</td>
</tr>
<tr>
<td>AVX-512</td>
<td>3857</td>
</tr>
<tr>
<td>FMA</td>
<td>32</td>
</tr>
<tr>
<td>KNC</td>
<td>601</td>
</tr>
<tr>
<td>SVML</td>
<td>406</td>
</tr>
</tbody>
</table>

2019
Visual Conventions We Will Use

Memory

Increasing address

Registers

- **Commonly:**

```
    R3  R2  R1  R0
```

- **We will use**

```
    R0  R1  R2  R3
```

AVX Intrinsics (Focus Floating Point)

Data types

```c
    __m256 f; // = {float f0, f1, f2, f3, f4, f5, f6, f7}
    __m256d d; // = {double d0, d1, d3, d4}
    __m256i i; // 32 8-bit, 16 16-bit, 8 32-bit, or 4 64-bit
```
AVX Intrinsics (Focus Floating Point)

Instructions

- **Naming convention**: \_mm256\_<intrin_op>\_<suffix>
- **Example**:

```
// a is 32-byte aligned
double a[4] = {1.0, 2.0, 3.0, 4.0};
__m256d t = __mm256_load_pd(a);
```

- **Same result as**

```
__m256d t = __mm256_set_pd(4.0, 3.0, 2.0, 1.0)
```

AVX Intrinsics

Native instructions (one-to-one with assembly)

- \_mm256\_load\_pd() ↔ vmovapd
- \_mm256\_add\_pd() ↔ vaddpd
- \_mm256\_mul\_pd() ↔ vmulpd

Multi instructions (map to several assembly instructions)

- \_mm256\_set\_pd()
- \_mm256\_set1\_pd()

Macros and helpers

- \_MM\_SHUFFLE()
Intel Intrinsics Guide

Great resource to quickly find the right intrinsics

Has latency and throughput information for many instructions

*Note:* Intel measures throughput in cycles, i.e., really shows $1/\text{throughput}$. Example: Intel throughput 0.33 means throughput is 3 ops/cycle.

What Are the Main Issues?

Alignment is important (256 bit = 32 byte)

You need to code explicit loads and stores

Overhead through shuffles

Not all instructions in SSE (AVX) have a counterpart in AVX (or AVX-512)

*Reason:* building in hardware an AVX unit by pasting together 2 SSE units is easy (e.g., vaddpd is just 2 parallel addpd); if SSE “lanes” need to be crossed it is expensive
### SSE vs. AVX vs. AVX-512

<table>
<thead>
<tr>
<th></th>
<th>SSE</th>
<th>AVX</th>
<th>AVX-512</th>
</tr>
</thead>
<tbody>
<tr>
<td>float, double</td>
<td>4-way, 2-way</td>
<td>8-way, 4-way</td>
<td>16-way, 8-way</td>
</tr>
<tr>
<td>register</td>
<td>16 x 128 bits: %xmm0 - %xmm15</td>
<td>16 x 256 bits: %ymm0 - %ymm15</td>
<td>32 x 512 bits: %zmm0 - %zmm31</td>
</tr>
<tr>
<td></td>
<td>The lower halves are the %xms</td>
<td>The lower halves are the %yms</td>
<td></td>
</tr>
<tr>
<td>assembly ops</td>
<td>addps, mulpd, ...</td>
<td>vaddps, vmulpd</td>
<td>vaddps, vmulpd</td>
</tr>
<tr>
<td>intrinsics data type</td>
<td>__m128, __m128d</td>
<td>__m256, __m256d</td>
<td>__m512, __m512d</td>
</tr>
<tr>
<td>intrinsics instructions</td>
<td>_mm_load_ps, _mm_add_ps, ...</td>
<td>_mm256_load_ps, _mm256_add_ps</td>
<td>_mm512_load_ps, _mm512_add_ps</td>
</tr>
</tbody>
</table>

*Mixing SSE and AVX may incur penalties*

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### AVX Intrinsics

- Load and store
- Constants
- Arithmetic
- Comparison
- Conversion
- Shuffles
Load and Stores

### Loads and Stores

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<thead>
<tr>
<th>Intrinsic Name</th>
<th>Operation</th>
<th>Corresponding AVX Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_load_pd</td>
<td>Load four double values, address aligned</td>
<td>VMOVAPD ymm, mem</td>
</tr>
<tr>
<td>_mm256_loadu_pd</td>
<td>Load four double values, address unaligned</td>
<td>VMOVPD ymm, mem</td>
</tr>
<tr>
<td>_mm256_maskload_pd</td>
<td>Load four double values using mask</td>
<td>VMASKMOVPD ymm, mem</td>
</tr>
<tr>
<td>_mm256_broadcast_sd</td>
<td>Load one double value into all four words</td>
<td>VBROADCASTSD ymm, mem</td>
</tr>
<tr>
<td>_mm256_broadcast_pd</td>
<td>Load a pair of double values into the lower and higher part of vector.</td>
<td>VBROADCASTSD ymm, mem</td>
</tr>
<tr>
<td>_mm256_loadp</td>
<td>Load double values from memory using indices.</td>
<td>VGATHERPD ymm, mem, ymm</td>
</tr>
<tr>
<td>_mm256_set1_pd</td>
<td>Set all four words with the same value</td>
<td>Composite</td>
</tr>
<tr>
<td>_mm256_set_pd</td>
<td>Set four values</td>
<td>Composite</td>
</tr>
<tr>
<td>_mm256_setr_pd</td>
<td>Set four values, in reverse order</td>
<td>Composite</td>
</tr>
<tr>
<td>_mm256_setzero_pd</td>
<td>Clear all four values</td>
<td>VxorPD</td>
</tr>
<tr>
<td>_mm256_set_m128d</td>
<td>Set lower and higher 128-bit parts</td>
<td>VINSERTF128</td>
</tr>
</tbody>
</table>

Tables show only most important instructions in category

May incur a significant performance penalty

Load and Stores

```
1.0 2.0 3.0 4.0

p

1.0 2.0 3.0 4.0

a

a = _mm256_load_pd(p); // p 32-byte aligned
a = _mm256_loadu_pd(p); // p not aligned
```

Load pd on unaligned pointer: seg fault

May incur a significant performance penalty

Load and Stores

```skylake
Lat = 7
Tp = 2
```
### Loads and Stores

**Skylake:**
- Lat = -
- Tp = -

#### 1.0

```c
a = _mm256_broadcast_pd(p1); // p1 any alignment
```

#### 2.0

```c
b = _mm256_broadcast_sd(p2); // p2 any alignment
```

#### 3.0

**Skylake:**
- Lat = -
- Tp = -

#### 1.0

```c
a = _mm256_maskload_pd(p, mask); // p any alignment
```

- mask = 
  - 0x0
  - 0xff

#### 2.0

```c
_mm256_load_pd(p)
```
Loads and Stores

```
a = _mm256_i64gather_pd(p, offset, 8); // p any alignment
```

scale = \{1,2,4,8\}
above: scale = 8 = size of double

Stores Analogous to Loads

<table>
<thead>
<tr>
<th>Intrinsic Name</th>
<th>Operation</th>
<th>Corresponding AVX instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_store_pd</td>
<td>Store four values, address aligned</td>
<td>VMOVAPD</td>
</tr>
<tr>
<td>_mm256_storeu_pd</td>
<td>Store four values, address unaligned</td>
<td>VMOVUPD</td>
</tr>
<tr>
<td>_mm256_maskstore_pd</td>
<td>Store four values using mask</td>
<td>VMASKMOVDP</td>
</tr>
<tr>
<td>_mm256_storeu2_m128d</td>
<td>Store lower and higher 128-bit parts into</td>
<td>Composite</td>
</tr>
<tr>
<td></td>
<td>different memory locations</td>
<td></td>
</tr>
<tr>
<td>_mm256_stream_pd</td>
<td>Store values without caching, address aligned</td>
<td>VMOVNTPD</td>
</tr>
</tbody>
</table>
### Constants

<table>
<thead>
<tr>
<th>LSB</th>
<th>a</th>
<th>b</th>
<th>c</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.0</td>
<td>2.0</td>
<td>3.0</td>
<td>4.0</td>
</tr>
<tr>
<td>1.0</td>
<td>1.0</td>
<td>1.0</td>
<td>1.0</td>
</tr>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
</tbody>
</table>

\[
a = \_\text{mm256\_set\_pd}(4.0, 3.0, 2.0, 1.0);
\]

\[
b = \_\text{mm256\_set1\_pd}(1.0);
\]

\[
c = \_\text{mm256\_setzero\_pd};
\]

### Arithmetic

<table>
<thead>
<tr>
<th>Intrinsic Name</th>
<th>Operation</th>
<th>Corresponding AVX Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_add_pd</td>
<td>Addition</td>
<td>VADDPD</td>
</tr>
<tr>
<td>_mm256_sub_pd</td>
<td>Subtraction</td>
<td>VSUBPD</td>
</tr>
<tr>
<td>_mm256_addsub_pd</td>
<td>Alternatively add and subtract</td>
<td>VADDSUBPD</td>
</tr>
<tr>
<td>_mm256_hadd_pd</td>
<td>Half addition</td>
<td>VHADDPD</td>
</tr>
<tr>
<td>_mm256_hsub_pd</td>
<td>Half subtraction</td>
<td>VHSUBPD</td>
</tr>
<tr>
<td>_mm256_mul_pd</td>
<td>Multiplication</td>
<td>VMULPD</td>
</tr>
<tr>
<td>_mm256_div_pd</td>
<td>Division</td>
<td>VDIVPD</td>
</tr>
<tr>
<td>_mm256_sqrt_pd</td>
<td>Squared Root</td>
<td>VSQRTPD</td>
</tr>
<tr>
<td>_mm256_max_pd</td>
<td>Computes Maximum</td>
<td>VMAXPD</td>
</tr>
<tr>
<td>_mm256_min_pd</td>
<td>Computes Minimum</td>
<td>VMINPD</td>
</tr>
<tr>
<td>_mm256_ceil_pd</td>
<td>Computes Ceiling</td>
<td>VROUNDPD</td>
</tr>
<tr>
<td>_mm256_floor_pd</td>
<td>Computes Floor</td>
<td>VROUNDPD</td>
</tr>
<tr>
<td>_mm256_round_pd</td>
<td>Round</td>
<td>VROUNDPD</td>
</tr>
<tr>
<td>_mm256_dp_ps</td>
<td>Single precision dot product</td>
<td>VDPPS</td>
</tr>
<tr>
<td>_mm256_fmadd_pd</td>
<td>Fused multiply-add</td>
<td>VFMAADD132pd</td>
</tr>
<tr>
<td>_mm256_fmsub_pd</td>
<td>Fused multiply-subtract</td>
<td>VFMSUB132pd</td>
</tr>
<tr>
<td>_mm256_fmadadd_pd</td>
<td>Alternatively fmadd, fmsub</td>
<td>VFMAADDSUB132pd</td>
</tr>
</tbody>
</table>

Tables show only most important instructions in category
Arithmetic

\[
\begin{array}{cccc}
1.0 & 2.0 & 3.0 & 4.0 \\
0.5 & 1.5 & 2.5 & 3.5 \\
1.5 & 3.5 & 5.5 & 7.5 \\
\end{array}
\]

\[c = \_mm256\_add\_pd(a, b);\]

\textit{analogous:}

\[c = \_mm256\_sub\_pd(a, b);\]

\[c = \_mm256\_mul\_pd(a, b);\]

Example

\[
\text{void addindex(double } *x, \text{ int } n) \{
\text{for (int } i = 0; i < n; i++)
\text{x[i] = x[i] + i;}
\}
\]

Vectorization by drawing:

\[
\text{load\_pd} \rightarrow \text{add\_pd} \rightarrow \text{store\_pd}
\]
Example

```c
void addindex(double *x, int n) {
    for (int i = 0; i < n; i++)
        x[i] = x[i] + 1;
}
```

```c
#include <immintrin.h>

// n a multiple of 4, x is 32-byte aligned
void addindex_vec1(double *x, int n) {
    __m256d index, x_vec;
    for (int i = 0; i < n; i += 4) {
        x_vec = _mm256_load_pd(x+i);
        index = _mm256_set_pd(i+3, i+2, i+1, i); // create vector with indexes
        x_vec = _mm256_add_pd(x_vec, index); // add the two
        _mm256_store_pd(x+i, x_vec); // store back
    }
}
```

Is this the best solution?

No! _mm256_set_pd may be too expensive

Example

```c
void addindex(double *x, int n) {
    for (int i = 0; i < n; i++)
        x[i] = x[i] + i;
}
```

```c
#include <immintrin.h>

// n a multiple of 4, x is 32-byte aligned
void addindex_vec2(double *x, int n) {
    __m256d x_vec, init, incr, ind;
    ind = _mm256_set_pd(3, 2, 1, 0); // load 4 doubles
    incr = _mm256_set1_pd(4); // create vector with indexes
    for (int i = 0; i < n; i += 4) {
        x_vec = _mm256_load_pd(x+i); // load 4 doubles
        x_vec = _mm256_add_pd(x_vec, ind); // add the two
        ind = _mm256_add_pd(ind, incr); // update ind
        _mm256_store_pd(x+i, x_vec); // store back
    }
}
```

Code style helps with performance! Why?
Arithmetic

```
c = _mm256_max_pd(a, b);
```

Arithmetic

```
c = _mm256_addsub_pd(a, b);
```
### Arithmetic

- **a** = 1.0, 2.0, 3.0, 4.0
- **b** = 0.5, 1.5, 2.5, 3.5
- **c** = \( \text{mm256\_hadd\_pd}(a, b) \);

**analogous:**

- **c** = \( \text{mm256\_hsub\_pd}(a, b) \);

*Does not cross between 128-bit lanes*

### Example

// n is even, low pass filter on complex numbers
// output z is in interleaved format

```c
void clp(double *re, double *im, double *z, int n) {
    for (int i = 0; i < n; i+=2) {
        z[i] = (re[i] + re[i+1])/2;
        z[i+1] = (im[i] + im[i+1])/2;
    }
}
```

**Skylake:**

Lat = -
Tp = -
Example

```c
#include <immintrin.h>

// n a multiple of 4, re, im, z are 32-byte aligned
void clp_vec(double *re, double *im, double *z, int n) {
    __m256d half, v1, v2, avg;
    half = _mm256_set1_pd(0.5); // set vector to all 0.5
    for(int i = 0; i < n; i+=4) {
        v1 = _mm256_load_pd(re+i); // load 4 doubles of re
        v2 = _mm256_load_pd(im+i); // load 4 doubles of im
        avg = _mm256_hadd_pd(v1, v2); // add pairs of doubles
        avg = _mm256_mul_pd(avg, half); // multiply with 0.5
        _mm256_store_pd(z+i, avg); // save result
    }
}
```

// n is even, low pass filter on complex numbers
// output z is in interleaved format
void clp(double *re, double *im, double *z, int n) {
    for(int i = 0; i < n; i+=2) {
        z[i] = (re[i] + re[i+1])/2;
        z[i+1] = (im[i] + im[i+1])/2;
    }
}
```

Arithmetic (FMA)

```
LSB
\[ \begin{array}{cccc}
1.0 & 2.0 & 3.0 & 4.0 \\
\end{array} \]

LSB
\[ \begin{array}{cccc}
0.5 & 0.5 & 0.5 & 0.5 \\
\end{array} \]

LSB
\[ \begin{array}{cccc}
0.5 & 1.5 & 2.5 & 3.5 \\
\end{array} \]

LSB
\[ \begin{array}{cccc}
1.0 & 2.5 & 4.0 & 5.5 \\
\end{array} \]
```

**Skylake:**
Lat = 4
Tp = 2
```
d = _mm256_fmadd_pd(a, b, c);

analogous:
d = _mm256_fmsub_pd(a, b, c);

scalar FMA (128 bit):
d = _mm_fmadd_sd(a, b, c);
```
Example

```c
#include <immintrin.h>

void complex_square_fma(double *a, double *x, double *y, int n) {
    __m128d re, im, a_re, a_im, two;
    two = _mm_set_sd(2.0);
    a_re = __mm_set_sd(a[0]);
    a_im = __mm_set_sd(a[1]);
    for (int i = 0; i < n; i+=2) {
        x_re = __mm_load_sd(x+i);
        x_im = __mm_load_sd(x+i+1);
        re = __mm_fmadd_sd(x_re, x_re, a_re);
        im = __mm_mul_sd(two, x_re);
        im = __mm_fmadd_sd(im, x_im, a_im);
        _mm_store_sd(y+i, re);
        _mm_store_sd(y+i+1, im);
    }
}
```

// y = a + x^2 on complex numbers, a is constant
void complex_square(double *a, double *x, double *y, int n) {
    for (int i = 0; i < n; i+=2) {
        y[i] = a[0] + x[i]*x[i] - x[i+1]*x[i+1];
        y[i+1] = a[1] + 2.0*x[i]*x[i+1];
    }
}

Arithmetic

```c
__m256 __m256_dp_ps(__m256 a, __m256 b, const int mask)
```

Computes the pointwise product of a and b and writes a selected sum of the resulting numbers into selected elements of c; the others are set to zero. The selections are encoded in the mask. (Only for floats)

**Example:** mask = 117 = 01110101

```
LSB 1.0 2.0 3.0 4.0 0.5 1.5 2.5 3.5 a (low half) b (low half)

0.5 3.0 7.5 14.0

LSB 11.0 0 11.0 0 c (low half)
```

Same is done for the upper half
### Comparisons

<table>
<thead>
<tr>
<th>Intrinsic Name</th>
<th>Macro for operation</th>
<th>Operation</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_cmp_pd</td>
<td>_CMP_EQ_OQ</td>
<td>Equal</td>
</tr>
<tr>
<td>(VCMPPD)</td>
<td>_CMP_EQ_UQ</td>
<td>Equal (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_GE_OQ</td>
<td>Greater than or equal</td>
</tr>
<tr>
<td></td>
<td>_CMP_GT_OQ</td>
<td>Greater than</td>
</tr>
<tr>
<td></td>
<td>_CMP_LE_OQ</td>
<td>Less than or equal</td>
</tr>
<tr>
<td></td>
<td>_CMP_LT_OQ</td>
<td>Less than</td>
</tr>
<tr>
<td></td>
<td>_CMP_NEQ_OQ</td>
<td>Not equal</td>
</tr>
<tr>
<td></td>
<td>_CMP_NEQ_UQ</td>
<td>Not Equal (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_NGE_UQ</td>
<td>Not Greater Than or Equal (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_NGT_UQ</td>
<td>Not Greater Than (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_NLE_UQ</td>
<td>Not Less Than or Equal (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_NLT_UQ</td>
<td>Not Less Than (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_TRUE_UQ</td>
<td>True (unordered)</td>
</tr>
<tr>
<td></td>
<td>_CMP_FALSE_OQ</td>
<td>False</td>
</tr>
<tr>
<td></td>
<td>_CMP_ORD_Q</td>
<td>Ordered</td>
</tr>
<tr>
<td></td>
<td>_CMP_UNORD_Q</td>
<td>Unordered</td>
</tr>
</tbody>
</table>

Tables show only most important instructions in category

---

### Comparisons

```
1.0  2.0  3.0  4.0  a
-----
1.0  1.5  3.0  3.5  b

c = _mm256_cmp_pd(a, b, _CMP_EQ_OQ);
```

Each field:
- `0xfff...f` if true
- `0x0` if false

Return type: __m256d

**analogous:**
- 
- 

```
c = _mm256_cmp_pd(a, b, _CMP_GE_OQ);
c = _mm256_cmp_pd(a, b, _CMP_LT_OQ);
```
Example

```c
void fcond(double *x, size_t n) {
    int i;
    for(i = 0; i < n; i++) {
        if(x[i] > 0.5)
            x[i] += 1.1;
        else
            x[i] -= 1.1;
    }
}
```

Example

```c
#include <xmmintrin.h>
void fcond_vec1(double *x, size_t n) {
    int i;
    __m256d vt, vmask, vp, vm, vr, ones, mones, thresholds;
    ones   = _mm256_set1_pd(1);
    mones  = _mm256_set1_pd(-1);
    thresholds = _mm256_set1_pd(0.5);
    for(i = 0; i < n; i+=4) {
        vt = _mm256_load_pd(x+i);
        vmask = _mm256_cmp_pd(vt, thresholds, _CMP_GT_OQ);
        vp   = _mm256_and_pd(vmask, ones);
        vm   = _mm256_andnot_pd(vmask, mones);
        vr   = _mm256_add_pd(vt, _mm256_or_pd(vp, vm));
        _mm256_store_pd(x+i, vr);
    }
}
## Vectorization

\[ \text{Vectorization} = \]

![Image](www.druckundbestell.de)

---

### Conversion

<table>
<thead>
<tr>
<th>Intrinsic Name</th>
<th>Operation</th>
<th>Corresponding AVX Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_cvtepi32_pd</td>
<td>Convert from 32-bit integer</td>
<td>VCVTQ2PD</td>
</tr>
<tr>
<td>_mm256_cvtepi32_ps</td>
<td>Convert from 32-bit integer</td>
<td>VCVTQ2PS</td>
</tr>
<tr>
<td>_mm256_cvtpd_epi32</td>
<td>Convert to 32-bit integer</td>
<td>VCVTQ2DQ</td>
</tr>
<tr>
<td>_mm256_cvtps_epi32</td>
<td>Convert to 32-bit integer</td>
<td>VCVTQ2DQ</td>
</tr>
<tr>
<td>_mm256_cvtps_pd</td>
<td>Convert from floats</td>
<td>VCVTQ2PD</td>
</tr>
<tr>
<td>_mm256_cvtpd_ps</td>
<td>Convert to floats</td>
<td>VCVTQ2PS</td>
</tr>
<tr>
<td>_mm256_cvtpd_epi32</td>
<td>Convert to 32-bit integer with truncation</td>
<td>VCVTQ2DQ</td>
</tr>
<tr>
<td>_mm256_cvtsd_f64</td>
<td>Extract</td>
<td>MOVSD</td>
</tr>
<tr>
<td>_mm256_cvtsd_f32</td>
<td>Extract</td>
<td>MOVSS</td>
</tr>
</tbody>
</table>

Tables show only most important instructions in category
Conversion

```c
double _mm256_cvtsd_f64(__m256d a)
```

```
| 1.0 | 2.0 | 3.0 | 4.0 |
```

```c
double d;
d = _mm_cvtsd_f64(a);
```

---

Conversion

```c
__m256d _mm256_cvtepi32_pd(__m128i a)
```

```
convert
```

```c
int s INT8
```

```
doubles
```

See also:

```c
__m256d _mm256_cvtepi64_pd(__m256i a)
```

```c
__m256d _mm256_cvtepu32_pd(__m256i a)
```

```c
__m256d _mm256_cvtepu64_pd(__m256i a)
```
Shuffles

<table>
<thead>
<tr>
<th>Intrinsic Name</th>
<th>Operation</th>
<th>Corresponding AVX Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>_mm256_unpackhi_pd</td>
<td>Unpack High</td>
<td>VUNPCKHPD</td>
</tr>
<tr>
<td>_mm256_unpacklo_pd</td>
<td>Unpack Low</td>
<td>VUNPCKLDP</td>
</tr>
<tr>
<td>_mm256_movemask_pd</td>
<td>Create four-bit mask</td>
<td>VMOVMSKPD</td>
</tr>
<tr>
<td>_mm256_movedup_pd</td>
<td>Duplicates</td>
<td>VMOVDDUP</td>
</tr>
<tr>
<td>_mm256_blend_pd</td>
<td>Selects data from 2 sources using constant mask</td>
<td>VBLENDPD</td>
</tr>
<tr>
<td>_mm256_blendv_pd</td>
<td>Selects data from 2 sources using variable mask</td>
<td>VBLENDVDPD</td>
</tr>
<tr>
<td>_mm256_insertf128_pd</td>
<td>Insert 128-bit value into packed array elements selected by index.</td>
<td>VINSETFT128</td>
</tr>
<tr>
<td>_mm256_extractf128_pd</td>
<td>Extract 128-bits selected by index.</td>
<td>VEXTRACTFT128</td>
</tr>
<tr>
<td>_mm256_shuffle_pd</td>
<td>Shuffle</td>
<td>VSHUFPD</td>
</tr>
<tr>
<td>_mm256_permute_pd</td>
<td>Permute</td>
<td>VPERMILPD</td>
</tr>
<tr>
<td>_mm256_permute4x64_pd</td>
<td>Permute 64-bits elements</td>
<td>VPERMPD</td>
</tr>
<tr>
<td>_mm256_permute2f128_pd</td>
<td>Permute 128-bits elements</td>
<td>VPERM2F128</td>
</tr>
</tbody>
</table>

Tables show only most important instructions in category

Skylake:  
Lat = 1  
Tp = 1

Does not cross between 128-bit lanes
Shuffles

```c
__m256d_mm256_blendv_pd(__m256d a, __m256d b, __m256d mask)
```

Result is filled in each position by an element of a or b in the same position as specified by mask

**Example:**

```plaintext
LSB
0x0 0xffffffff 0x0 0x0 mask

<table>
<thead>
<tr>
<th></th>
<th>1.0</th>
<th>2.0</th>
<th>3.0</th>
<th>4.0</th>
</tr>
</thead>
<tbody>
<tr>
<td>LSB</td>
<td>a</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th></th>
<th>0.5</th>
<th>1.5</th>
<th>2.5</th>
<th>3.5</th>
</tr>
</thead>
<tbody>
<tr>
<td>LSB</td>
<td>b</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th></th>
<th>1.0</th>
<th>1.5</th>
<th>3.0</th>
<th>4.0</th>
</tr>
</thead>
<tbody>
<tr>
<td>LSB</td>
<td>c</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
```

see also _mm256_blend_pd:
same with integer mask, \( T_p = 3! \)

---

**Example (Continued From Before)**

```c
void fcond(double *x, size_t n) {
    int i;
    for(i = 0; i < n; i++) {
        if(x[i] > 0.5)
            x[i] += 1.0;
        else
            x[i] -= 1.0;
    }
}
```

```c
#include <immintrin.h>
void fcond_vec2(double *x, size_t n) {
    __m256d vt, vmask, vp, vn, vr, ones, mones, thresholds;
    ones       = _mm256_set1_pd(1.0);
mones      = _mm256_set1_pd(-1.0);
    thresholds = _mm256_set1_pd(0.5);
    for(i = 0; i < n; i+=4) {
        vt = _mm256_load_pd(x+i);
        vmask = _mm256_cmp_pd(vt, thresholds, _CMP_GT_OQ);
        vb = _mm256_blendv_pd(mones, ones, vmask);
        vr = _mm256_add_pd(vt, vb);
        _mm256_store_pd(x+i, vr);
    }
}
```

Coffee Lake, gcc 7.3.0
-O3 -march=skylake -mfma

4–10x speedup
Example: Loading 4 Real Numbers from Arbitrary Memory Locations

Assumes all values are within one array

![Diagram showing loading process]

7 instructions, this is one way of doing it

Code For Previous Slide

```c
#include <immintrin.h>

__m256d LoadArbitrary(double *p0, double *p1, double *p2, double *p3) {
    __m256d a, b, c, d, e, f;
    a = __mm256_loadu_pd(p0);
    b = __mm256_loadu_pd(p1);
    c = __mm256_loadu_pd(p2-2);
    d = __mm256_loadu_pd(p3-2);
    e = __mm256_unpacklo_pd(a, b);
    f = __mm256_unpacklo_pd(c, d);
    return __mm256_blend_pd(e, f, 0b1100);
}
```

Example compilation:

```c
vmovupd ymm0, [rdi]
vmovupd ymm1, [-16+rdx]
vunpcklpd ymm2, ymm0, [rsi]
vunpcklpd ymm3, ymm1, [-16+rcx]
vblendpd ymm0, ymm2, ymm3, 12
```

No intrinsic for this instruction (Nov 2019)
Example: Loading 4 Real Numbers from Arbitrary Memory Locations (cont’d)

Whenever possible avoid the previous situation

Restructure algorithm and use the aligned
_mm256_load_pd()
Example: Loading 4 Real Numbers from Arbitrary Memory Locations (cont’d)

Example compilation:
```
vmovsd  xmm0  [rdi]
vmovsd  xmm1,  [rdx]
vmovhpd  xmm2,  xmm0,  [rsi]  // SSE register xmm2 written
vmovhpd  xmm3,  xmm1,  [rcx]
vinsertf128  ymm0,  ymm2,  xmm3,  1  // accessed as ymm2
```

Written in intrinsics (reverse-engineered):
```
#include <immintrin.h>

__m256d myArbitraryLoad2(double *a, double *b, double *c, double *d) {
    _m128d t1, t2, t3, t4;
    _m256d t5;
    t1 = _mm_load_sd(a);  // SSE
    t2 = _mm_loadh_pd(t1, b);  // SSE
    t3 = _mm_load_sd(c);  // SSE
    t4 = _mm_loadh_pd(t3, d);  // SSE
    t5 = _mm256_castpd128_pd256(t2);  // cast _m128d -> _m256d
    return _mm256_insertf128_pd(t5, t4, 1);
}
```

Example: Loading 4 Real Numbers from Arbitrary Memory Locations

Picture for previous slide (this solution always works):
Example: Loading 4 Real Numbers from Arbitrary Memory Locations (cont’d)

Do not do this (why?):

```c
__declspec(align(32)) double g[4];
__m256d vf;
g[0] = *p0;
g[1] = *p1;
g[2] = *p2;
g[3] = *p3;
vf = _mm256_load_pd(g);
```

Shuffles

```
__m256d _mm256_shuffle_pd(__m256d a, __m256d b, const int mask)
```

Skylake:
Lat = 1
Tp = 1

Does not cross between 128-bit lanes
Shuffles

\texttt{__m256d \_mm256_permute\_pd(__m256d a, int mask)}

Shuffle elements within 128-bits lanes.

Example:

\begin{itemize}
\item \text{LSB}:
  \begin{itemize}
  \item 1
  \item 2
  \item 3
  \item 4
  \end{itemize}
\item \text{a}
\item \text{1101 mask}
\item \text{LSB}:
  \begin{itemize}
  \item 2
  \item 1
  \item 4
  \item 4
  \end{itemize}
\item \text{c}
\end{itemize}

\textit{Does not cross between 128-bit lanes}

Shuffles

\texttt{__m256d \_mm256_permute4x64\_pd(__m256d a, int mask)}

Result is filled in each position by any element of a, as specified by mask

Example:

\begin{itemize}
\item \text{LSB}:
  \begin{itemize}
  \item 1
  \item 2
  \item 3
  \item 4
  \end{itemize}
\item \text{a}
\item \text{11010001 mask}
\item \text{LSB}:
  \begin{itemize}
  \item 2
  \item 1
  \item 2
  \item 4
  \end{itemize}
\item \text{c}
\end{itemize}

\textit{Somewhat more expensive due to shuffle between 128-bits lanes}
Vectorization With Intrinsics: Key Points

Use aligned loads and stores as much as possible

Minimize shuffle instructions

Minimize use of suboptimal arithmetic instructions.
E.g., \texttt{add\_pd} has higher throughput than \texttt{hadd\_pd}

Be aware of available instructions (\texttt{intrinsics guide!}) and their performance

SIMD Extensions and AVX

AVX intrinsics

\textit{Compiler vectorization}

References:
\texttt{Intel \textit{icc manual}} (look for auto vectorization)
Compiler Vectorization

Compiler flags
Aliasing
Proper code style
Alignment

How Do I Know the Compiler Vectorized?

vec-report

Look at assembly: vmulpd, vaddpd, xxxpd

Generate assembly with source code annotation:
  - Visual Studio + icc: /Fas
  - icc on Linux/Mac: -S
Example

unvectorized: /Qvec-

```c
void myadd(double *a, double *b, const int n) {
  for (int i = 0; i < n; i++)
    a[i] = a[i] + b[i];
}
```

```c
;;    a[i] = a[i] + b[i];
vmovsd xmm0, DWORD PTR [rcx+rax*4]
vaddsd xmm0, DWORD PTR [rdx+rax*4]
vmovsd DWORD PTR [rcx+rax*4], xmm0
```

vectorized:

```c
comore
;;    a[i] = a[i] + b[i];
vmovsd xmm0, DWORD PTR [rcx+r11*4]
vaddsd xmm0, DWORD PTR [rdx+r11*4]
vmovsd DWORD PTR [rcx+r11*4], xmm0

vmovupd ymm0, YMMWORD PTR [rdx+r10*4]
vmovupd ymm1, YMMWORD PTR [16+rdx+r10*4]
vaddpd ymm0, ymm0, YMMWORD PTR [rcx+r10*4]
vaddpd ymm1, ymm1, YMMWORD PTR [16+rcx+r10*4]
vmovupd YMMWORD PTR [rcx+r10*4], ymm0
vmovupd YMMWORD PTR [16+rcx+r10*4], ymm1
comore
```

Are These Programs Equivalent?

P1:
```c
for (i = 0; i < n; i++) // n even
  a[i] = a[i] + b[i];
```

P2:
```c
for (i = 0; i < n; i+=2) // n even
{
  s1 = a[i];
  s2 = a[i+1];
  t1 = b[i];
  t2 = b[i+1];
  s1 = s1 + t1;
  s2 = s2 + t2;
  a[i] = s1;
  a[i+1] = s2;
}
```

No! Possible aliasing
Aliasing

\[
\begin{align*}
\text{for } (i = 0; i < n; i++) \\
\quad a[i] &= a[i] + b[i];
\end{align*}
\]

Cannot be vectorized in a straightforward way due to potential aliasing.

However, in this case compiler can insert runtime check:

\[
\begin{align*}
\text{if } (a + n < b \ | \ b + n < a) \\
\quad /* \text{vectorized loop */} \\
\text{...} \\
\else \\
\quad /* \text{serial loop */} \\
\text{...}
\end{align*}
\]

Removing Aliasing

Globally with compiler flag:
- `-fno-alias`, `/Oa`
- `-fargument-noalias`, `/Qalias-args` - (function arguments only)

For one loop: pragma

```c
void add(double *a, double *b, int n) {
    #pragma ivdep
    for (i = 0; i < n; i++)
        a[i] = a[i] + b[i];
}
```

For specific arrays: restrict (needs compiler flag `-restrict`, `/Qrestrict`)

```c
void add(double *restrict a, double *restrict b, int n) {
    for (i = 0; i < n; i++)
        a[i] = a[i] + b[i];
}
```
Proper Code Style

Use countable loops = number of iterations known at runtime

- **Number of iterations is a:**
  - constant
  - loop invariant term
  - linear function of outermost loop indices

Countable or not?

```c
for (i = 0; i < n; i++)
a[i] = a[i] + b[i];
```

```c
void vsum(double *a, double *b, double *c) {
    int i = 0;

    while (a[i] > 0.0) {
        a[i] = b[i] * c[i];
        i++;
    }
}
```

Proper Code Style

Use arrays, structs of arrays, not arrays of structs

Ideally: unit stride access in innermost loop

```c
void mmm1(double *a, double *b, double *c) {
    int N = 100;
    int i, j, k;

    for (i = 0; i < N; i++)
        for (j = 0; j < N; j++)
            for (k = 0; k < N; k++)
                c[i*n + j] += a[i*n + k] * b[k*n + j];
}
```

```c
void mmm2(double *a, double *b, double *c) {
    int N = 100;
    int i, j, k;

    for (i = 0; i < N; i++)
        for (k = 0; k < N; k++)
            for (j = 0; j < N; j++)
                c[i*n + j] += a[i*n + k] * b[k*n + j];
}
```
Alignment

double *x = (double *) malloc(1024*sizeof(double));
int i;

for (i = 0; i < 1024; i++)
  x[i] = 1;

Without alignment information would require unaligned loads if vectorized. However, the compiler can peel the loop to start it at an aligned address: the generated assembly would mimic the below C code:

double *x = (double *) malloc(1024*sizeof(double));
int i;

peel = (unsigned long) x & 0x1f; /* x mod 32 */
if (peel != 0) {
  peel = (32 - peel)/sizeof(double);
  /* initial segment */
  for (i = 0; i < peel; i++)
    x[i] = 1;
}
/* 32-byte aligned access */
for (i = peel; i < 1024; i++)
  x[i] = 1;

Ensuring Alignment

Align arrays to 32-byte boundaries (see earlier discussion)

If compiler cannot analyze:

- Use pragma for loops

    double *x = (double *) malloc(1024*sizeof(double));
    int i;

    #pragma vector aligned
    for (i = 0; i < 1024; i++)
      x[i] = 1;

- For specific arrays:
  __assume_aligned(a, 32);

Use simple for loops. Avoid complex loop termination conditions – the upper iteration limit must be invariant within the loop. For the innermost loop in a nest of loops, you could set the upper limit iteration to be a function of the outer loop indices.

Write straight-line code. Avoid branches such as switch, goto, or return statements, most function calls, or if constructs that can not be treated as masked assignments.

Avoid dependencies between loop iterations or at the least, avoid read-after-write dependencies.

Try to use array notations instead of the use of pointers. C programs in particular impose very few restrictions on the use of pointers; aliased pointers may lead to unexpected dependencies. Without help, the compiler often cannot tell whether it is safe to vectorize code containing pointers.

Wherever possible, use the loop index directly in array subscripts instead of incrementing a separate counter for use as an array address.

Access memory efficiently:
- Favor inner loops with unit stride.
- Minimize indirect addressing.
- Align your data to 32 byte boundaries (for AVX instructions).

Choose a suitable data layout with care. Most multimedia extension instruction sets are rather sensitive to alignment.

Read the above website

```c
void myadd(double *a, double *b, const int n) {
    for (int i = 0; i < n; i++)
        a[i] = a[i] + b[i];
}
```

Assume:
- No aliasing information
- No alignment information

Can compiler vectorize?

In principle yes: through versioning

However, this causes code size blowup and is not feasible for large code
Compiler Vectorization

Understand the limitations

Carefully read the manual